

Hog Contest Rules

- Up to two people submit one entry;
 Max of one entry per person
- Slight rule changes
- Your score is the number of entries against which you win more than 50.00001% of the time
- Strategies are time-limited
- All strategies must be deterministic, pure functions of the players' scores
- All winning entries will receive extra credit
- The real prize: honor and glory
- See website for detailed rules

Fall 2011 Winners

Kaylee Mann Yan Duan & Ziming Li Brian Prike & Zhenghao Qian Parker Schuh & Robert Chatham

Fall 2012 Winners

Chenyang Yuan Joseph Hui

Fall 2013 Winners

Paul Bramsen Sam Kumar & Kangsik Lee Kevin Chen

Fall 2014 Winners

Alan Tong & Elaine Zhao Zhenyang Zhang Adam Robert Villaflor & Joany Gao Zhen Qin & Dian Chen Zizheng Tai & Yihe Li

Hog Contest Winners

Spring 2015 Winners

Sinho Chewi & Alexander Nguyen Tran Zhaoxi Li Stella Tao and Yao Ge

Fall 2015 Winners

Micah Carroll & Vasilis Oikonomou Matthew Wu Anthony Yeung and Alexander Dai

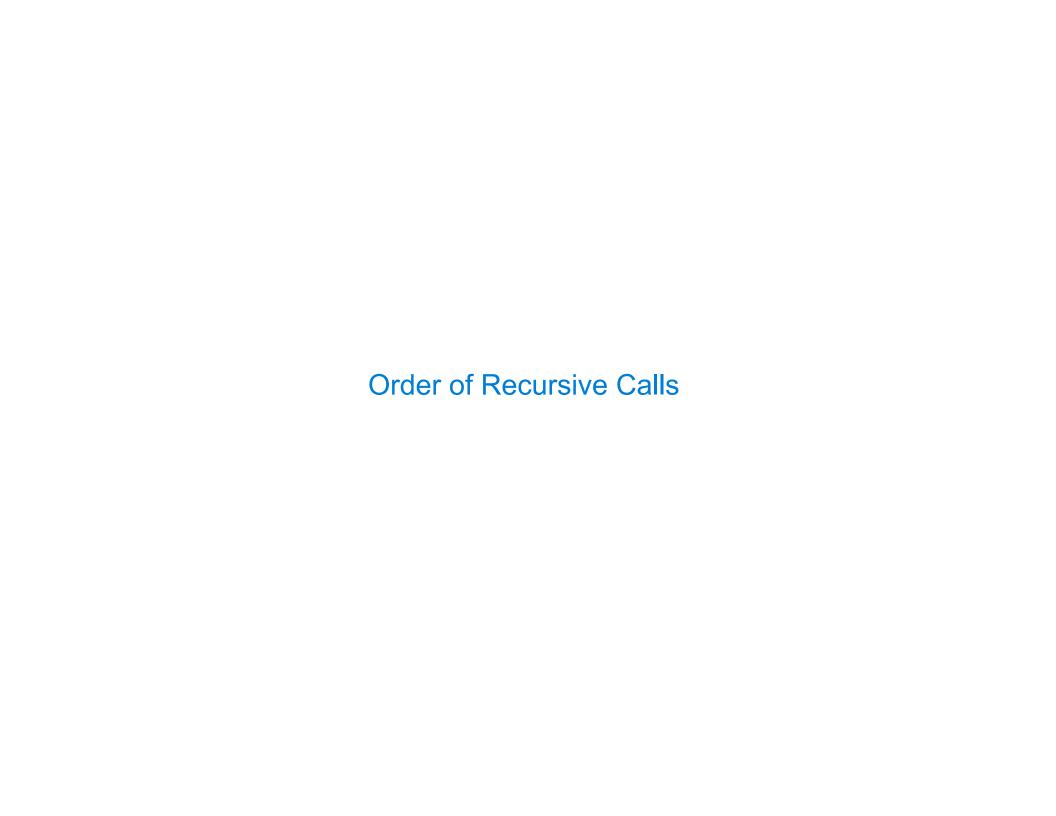
Spring 2016 Winners

Michael McDonald and Tianrui Chen Andrei Kassiantchouk Benjamin Krieges

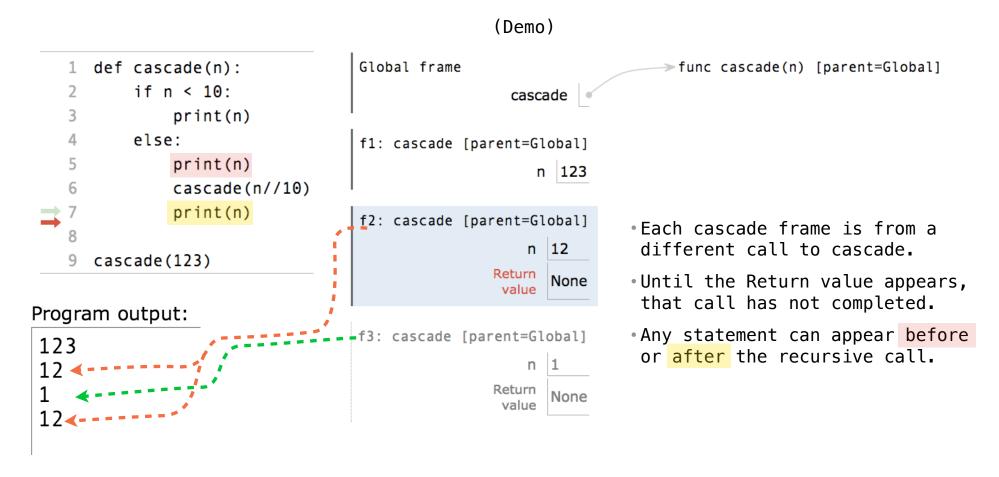
Spring 2017 Winners

Cindy Jin and Sunjoon Lee Anny Patino and Christian Vasquez Asana Choudhury and Jenna Wen Michelle Lee and Nicholas Chew

Fall 2017 Winners



The Cascade Function



Two Definitions of Cascade

(Demo)

```
def cascade(n):
    if n < 10:
        print(n)
        print(n)
        print(n)
        print(n)
        cascade(n//10)
        print(n)
        cascade(n//10)
        print(n)</pre>
```

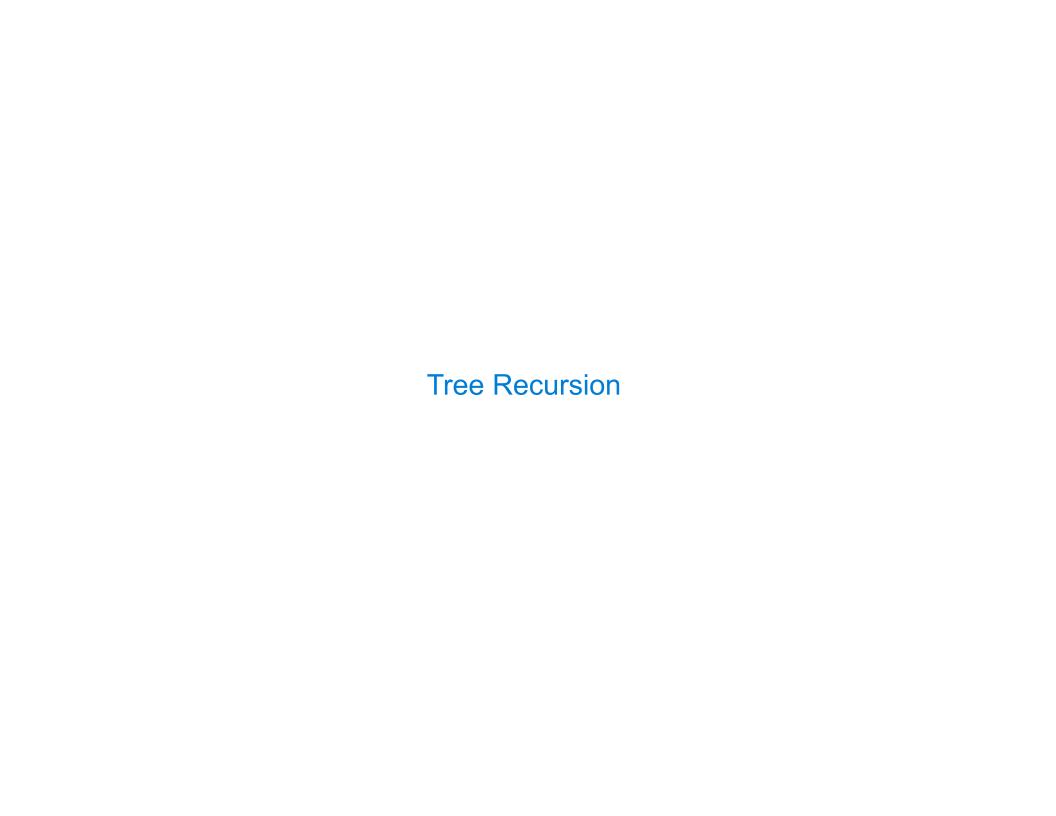
- If two implementations are equally clear, then shorter is usually better
- In this case, the longer implementation is more clear (at least to me)
- When learning to write recursive functions, put the base cases first
- Both are recursive functions, even though only the first has typical structure

Example: Inverse Cascade

Inverse Cascade

Write a function that prints an inverse cascade:

5



Tree Recursion

Tree—shaped processes arise whenever executing the body of a recursive function makes more than one recursive call

```
n: 0, 1, 2, 3, 4, 5, 6, 7, 8, ..., 35

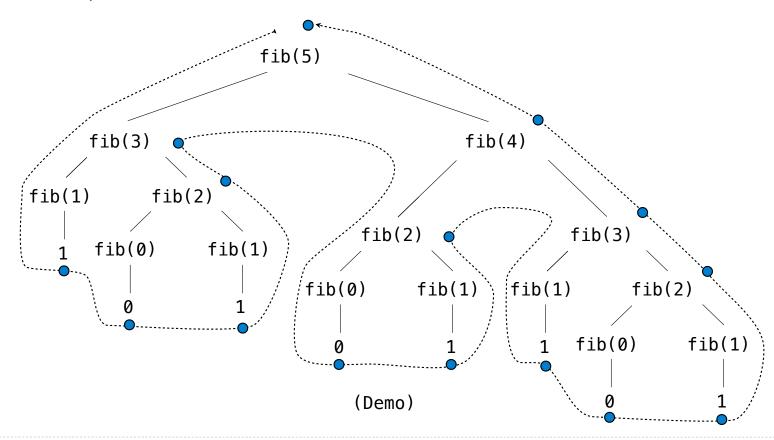
fib(n): 0, 1, 1, 2, 3, 5, 8, 13, 21, ..., 9,227,465
```

```
def fib(n):
    if n == 0:
        return 0
    elif n == 1:
        return 1
    else:
        return fib(n-2) + fib(n-1)
```



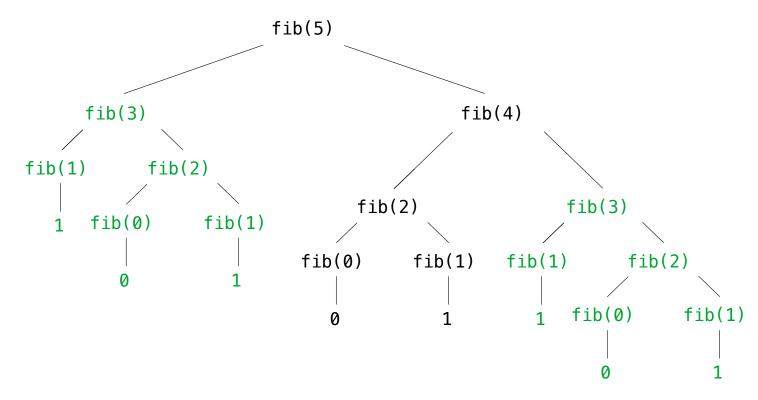
A Tree-Recursive Process

The computational process of fib evolves into a tree structure



Repetition in Tree-Recursive Computation

This process is highly repetitive; fib is called on the same argument multiple times



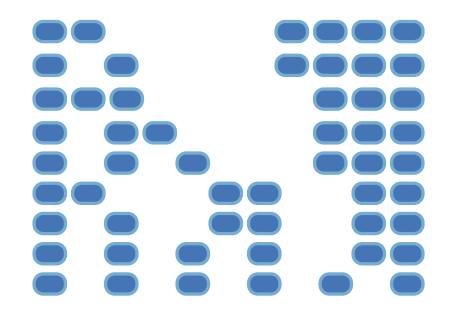
(We will speed up this computation dramatically in a few weeks by remembering results)

Example: Counting Partitions

Counting Partitions

The number of partitions of a positive integer n, using parts up to size m, is the number of ways in which n can be expressed as the sum of positive integer parts up to m in increasing order.

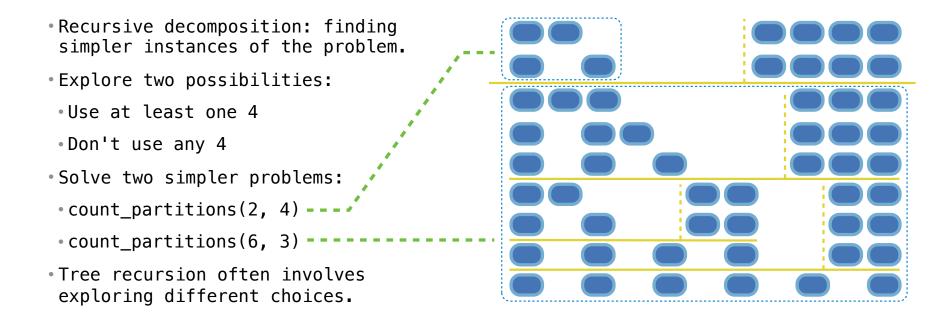
count_partitions(6, 4)



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```
def count_partitions(n, m):

    Recursive decomposition: finding

                                               if n == 0:
 simpler instances of the problem.
                                                   return 1
Explore two possibilities:
                                               elif n < 0:
                                                   return 0
•Use at least one 4
                                               elif m == 0:
• Don't use any 4
                                                   return 0
• Solve two simpler problems:
                                                else:
                                         \longrightarrow with m = count partitions(n-m, m)
• count partitions(2, 4) -----
                                                   without m = count partitions(n, m-1)
•count partitions(6, 3) -----
                                                    return with m + without m

    Tree recursion often involves

exploring different choices.
                                          (Demo)
```

<u>Interactive Diagram</u>